

- RYDE -

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Overview

RYDE is a digital signature scheme named after the **R**ank **S**Yndrome **D**Ecoding problem

- ◇ Fiat-Shamir (**FS**) based signature along with a Zero-Knowledge Proof of Knowledge (**PoK**)
- ◇ PoK built using the Multi-Party Computation in the Head (**MPCitH**) paradigm
- ◇ PoK relies on the hardness of the Rank Syndrome Decoding problem (**RSD**)

<https://pqc-ryde.org>

Agenda

- 1 Round 2 Updates
- 2 RSD Problem
- 3 Scheme Overview
- 4 Sizes & Performances
- 5 Advantages & Limitations

Round 2 Updates

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New results since Round 1

- ◇ New modeling for RSD [BFG⁺24]
- ◇ New MPCitH frameworks - **TCitH** [FR25] & **VOLeith** [BBD⁺23]

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- ◇ New modeling for RSD [BFG⁺24]
- ◇ New MPCitH frameworks - **TCitH** [FR25] & **VOLEitH** [BBD⁺23]

Modifications for Round 2

- ◇ v2.0.0 - Design update using the new modeling along with new MPCitH frameworks
- ◇ v2.0.1 - Implementation update
- ◇ v2.1.0 - Implementation update & Parameters fine-tuning

Round 2 Updates

RYDE Instance	Modeling	Proof System	Size (pk + sig.)
Round 1	Annihilator q -polynomial	MPCitH	6.1 - 7.6 kB
Round 2	Dual Support Decomposition	TCitH (& VOLEitH)	3.2 - 3.7 kB

Table 1: Modifications for RYDE (sizes are given for NIST-1 security level)

RSD Problem

Rank Metric

RYDE relies on **code-based cryptography** in the **rank metric** setting

$$\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{F}_{q^m}^n \quad \Leftrightarrow \quad \mathbf{M}_{\mathbf{x}} = \begin{bmatrix} x_{1,1} & \dots & x_{n,1} \\ \vdots & \ddots & \vdots \\ x_{1,m} & \dots & x_{n,m} \end{bmatrix} \in \mathbb{F}_q^{m \times n}$$

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- ◇ $\text{Supp}(\mathbf{x}) = \langle x_1, \dots, x_n \rangle_{\mathbb{F}_q}$
- ◇ $w_R(\mathbf{x}) = \text{rank}(\mathbf{M}_{\mathbf{x}})$

RSD Problem

Rank Syndrome Decoding Problem

Input

- Secret value $\mathbf{x} \in \mathbb{F}_{q^m}^n$ such that $w_R(\mathbf{x}) = r$
- Public values $\mathbf{H} \in \mathbb{F}_{q^m}^{(n-k) \times n}$ and $\mathbf{y} \in \mathbb{F}_{q^m}^{n-k}$ such that $\mathbf{xH}^\top = \mathbf{y}$

Goal

- Find $\tilde{\mathbf{x}} \in \mathbb{F}_{q^m}^n$ such that $\tilde{\mathbf{x}}\mathbf{H}^\top = \mathbf{y}$ and $w_R(\tilde{\mathbf{x}}) = r$

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RYDE relies on the hardness of the RSD problem (***without*** cyclic structure)

Scheme Overview

Modeling

RYDE v2 relies on the Dual Support Decomposition modeling for RSD [BFG⁺24]

- ◇ Natural modeling checking the weight of $\mathbf{x}$ using matrix decomposition
- ◇ Updated RSD parameter sets to minimize the witness size

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RYDE Instance	Modeling	Witness Size (for NIST-1 security level)	
Round 1	Annihilator q -polynomial	$[(r - 1)m + km] \cdot \log_2(q)$	93 B
Round 2	Dual Support Decomposition	$[(r - 1)m + r(n - r)] \cdot \log_2(q)$	45 B

Table 2: RYDE modeling and resulting witness sizes

Modeling

Protocol Overview

Public Input

- An instance $(\mathbf{H}, \mathbf{y})$ of the RSD problem

Private Input

- Coefficients $\mathbf{s}' \in \mathbb{F}_{q^m}^{(r-1)}$ of a basis $\mathbf{s} = (\mathbf{1} \ \mathbf{s}')$ of the support of $\mathbf{x}$
- Coefficients $\mathbf{C}' \in \mathbb{F}_q^{r \times (n-r)}$ of the support decomposition of $\mathbf{x}$ with respect to $\mathbf{s}$

Protocol

1. Verify the weight of $\mathbf{x}$ by computing $\mathbf{x} = \mathbf{s} \cdot (\mathbf{I}_r \ \mathbf{C}')$
2. Verify that $\mathbf{x}$ is a solution by checking $\mathbf{x}\mathbf{H}^\top = \mathbf{y}$

Proof System

MPCitH Frameworks

- ◇ Two recent improvements to the MPCitH paradigm - **TCitH** [FR25] & **VOLEitH** [BBD⁺23]
- ◇ TCitH and VOLEitH can be described using the PIOP formalism [Fen24]

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TCitH

- ◇ **5**-round protocol
- ◇ Computation over a small field
- ◇ Several protocol repetitions
- ◇ *Arguably* simpler

VOLEitH

- ◇ **7**-round protocol
- ◇ Computation over a large field
- ◇ One protocol execution
- ◇ Smaller signatures

Proof System

RYDE & TCitH vs VOLEitH

- ◇ TCitH and VOLEitH lead to comparable sizes for modeling with low multiplicative depth
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RYDE Instantiation

- ◇ RYDE is instantiated with the **TCitH** framework (with a VOLEitH variant also described)
- ◇ RYDE uses the *one tree* optimization for GGM trees [BBM⁺24]

Sizes & Performances

Implementation

Implementation Updates

- ◇ Overall improvement of the performances of the scheme
- ◇ Update of symmetric primitives (AES/Rijndael for some PRG, AES/Rijndael variant for cmt)
- ◇ Reported constant-time issues have been fixed [ABB⁺25]

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Fine-Tuning Parameters

- ◇ RSD parameters updated for NIST-5 security level based on performance considerations
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Benchmark & Ongoing Work

- ◇ Numbers reported for the fastest variant of the optimized implementation (avx2 & aes-ni)
- ◇ Ongoing work targeting additional performance improvements

Sizes & Performances

RYDE-1 Instance		sk	pk	sig.	Keygen	Sign	Verify
Round 1	Short	32 B	86 B	6.0 kB	33 K	23 M	20 M
Round 2 (v2.1.0)	Short	32 B	69 B	3.2 kB	34 K	18 M	15 M
Round 1	Fast	32 B	86 B	7.5 kB	33 K	5.4 M	4.4 M
Round 2 (v2.1.0)	Fast	32 B	69 B	3.6 kB	34 K	7.0 M	2.8 M
Round 2 (v2.1.0)	Faster	32 B	69 B	5.0 kB	34 K	1.7 M	0.9 M

Table 3: Sizes and performances (CPU cycles) of RYDE (TCitH) for NIST-1 security level

Sizes & Performances

RYDE-5 Instance		sk	pk	sig.	Keygen	Sign	Verify
Round 1	Short	64 B	188 B	22.9 kB	72 K	106 M	95 M
Round 2 (v2.1.0)	Short	64 B	133 B	12.7 kB	67 K	141 M	131 M
Round 1	Fast	64 B	188 B	29.2 kB	72 K	26 M	23 M
Round 2 (v2.1.0)	Fast	64 B	133 B	14.8 kB	67 K	29 M	24 M
Round 2 (v2.1.0)	Faster	64 B	133 B	20.9 kB	67 K	7.3 M	6.6 M

Table 4: Sizes and performances (CPU cycles) of RYDE (TCitH) for NIST-5 security level

Comparison to other schemes

- Stay tuned till the end of the session -

*Overview of MPCitH based Signatures using the **PQ-SORT** benchmarking tool*

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- ◇ *Size* - Small public keys & Competitive signature size
 $|\text{pk} + \text{sig.}| \Rightarrow$ **3.2 kB** for RYDE, **3.7 kB** for ML-DSA, **7.8 kB** for SLH-DSA (for NIST-1 level)

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Limitations

- ◇ *Size* - Quadratic growth of signature sizes with respect to security level
- ◇ *Performances* - Slower than lattice-based signature schemes
But competitive with many other post-quantum signatures

Thank you for your attention.

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